

Sandra Müller, Girona July 17, 2018

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Long games and Woodin cardinals

joint work with Juan Aguilera

We consider "long games". Let $A \subseteq {}^{\omega^\alpha} \omega \approx (\omega_\omega)^\alpha$ for some fixed $\alpha < \omega_1^\omega$.

I	n_0	n_2	...	n_w	n_{w+2}	...
II	n_1	n_3	...	n_{w+1}	n_{w+3}	...

Player I wins iff $(n_0, n_1, \dots) \in A$, o/w Player II wins.

Rmk: Determinacy for games of length $\omega \cdot (n+1)$ with \mathbb{II}^1 payoff implies determinacy for games of length ω with \mathbb{II}_{n+1}^1 payoff.

Idea: We can simulate projections by ω moves (one round) in a game, where we only consider the moves of one of the two players.

Example: Let $A = pB$, $B \subseteq (\omega^\omega)^2$, $B \in \mathbb{II}_1^1$.

G^*	I	x_0	x_2	...	y_0	y_1	y_2	...
	II	x_1	x_3	...	II plays something but we don't care what			

Player I wins iff $(x, y) \in B$

G	I	x_0	x_2	...	Player I wins iff $x \in A$.			
	II	x_1	x_3	...				

Let σ be a ws. for I in G^* . Then $\sigma \upharpoonright w$ (restricted to the first ω moves of the game) is a ws. for I in G .
Analogous for Player II.

Thm (Neeman): Let $\alpha > 1$ be a ctble ordinal and suppose [2] that there are $-1+\alpha$ Woodin cardinals with a # above them all. Then all games of length $\omega \cdot \alpha$ with II^1 payoff (in fact even $\omega^2 \cdot \text{II}^1$ payoff) are determined.

Note: Woodin showed this earlier for ordinals α of the form $\alpha = \omega \cdot \beta$, where β is additively closed.

Question: Is this result optimal, i.e. can we prove the converse?

Theorem: Suppose games of length $\omega \cdot (\omega + 1)$ with II^1 payoff are determined. Then there is a premouse with $\omega + 1$ Woodin cardinals

Rmk: In fact, we will consider a game of length $\omega \cdot \omega$ with II^1 payoff.

Note: The same proof works for games of length $\omega \cdot (\omega + n)$ with II^1 payoff and $\omega + n$ Woodin cardinals for every new.

Idea of the proof: From now on suppose games of length $\omega \cdot (\omega + 1)$ with II^1 payoff are determined.

We first show that there is a model of the form $M_1(A)$ for some $A \in P_{\omega_1}(\mathbb{R})$ s.t. $M_1(A) \upharpoonright \mathbb{R} = A$ and $M_1(A) \models \text{AD}$ (and in fact $M_1(A) \models \text{AD}^+$).

- Use AD to "generate" ω Woodin cardinals in a generic extension of $M_1(A)$ (with a Prikry-type forcing)
- Use a P-construction to add the remaining Woodin cardinal on top.

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Lemma 1: There is a club $C_1 \subseteq P_{\omega_1}(\text{TR})$ such that $\text{TR} \cap M_1(A) = A$ for all $A \in C_1$.

We thank John Steel for pointing out to us that a variant of our proof of Lemma 2 below shows Lemma 1.

Lemma 2: There is a club $C_2 \subseteq P_{\omega_1}(\text{TR})$ such that $M_1(A) \models \text{ZF} + \text{AD}$ for all $A \in C_2$.

Proof: Sps. not, i.e. there is a stationary set of $B \in P_{\omega_1}(\text{TR})$ with $M_1(B) \models \delta_B \not\models \text{AD}$.
↑ wdn. card in $M_1(B)$

Now we play the following game G_1 .

I	z	a	v_0, x_1	v_1, x_3	\dots	where
II		b	x_2	x_4	\dots	

- $z, x_1, x_2, \dots \in {}^{\omega}\omega$ (played as sequences of natural numbers)
- $a, b \in {}^{\omega}\omega$ are obtained by alternating moves of I and II
- $v_i \in \{0, 1\}$ are interpreted as truth values of formulae ϕ_i , where $\{\phi_i : i < \omega\}$ is a fixed enumeration of all $L_{\text{pm}}(\{x_i : i < \omega\})$ -formulae (st. x_i does not appear in ϕ_j if $j \leq i$).

→ Every play determines a complete theory T in the language $L_{\text{pm}}(\{x_i : i < \omega\})$.

Player I wins G_1 iff

$$(1) \quad x_1 \models_T a \oplus b$$

(2) For each $i < \omega$, T contains the sentence $x_i \in {}^{\omega}\omega$

and for each $j, m < \omega$, T contains the sentence $x_i(m) = j$ iff $x_i(m) = j$, where we let $x_0 = z$.

(3) Let m and n be fixed maps, mapping each $\lambda pm(fx_i : i < \omega^y)$ -formula ϕ to an even natural number m_ϕ (or n_ϕ) such that m and n are recursive, have disjoint ranges, and m_ϕ and n_ϕ are larger than $\max f_i \mid x_i \text{ occurs in } \phi$. [4]

For every formula $\phi(x)$ with one free variable, T contains the sentences

$$\exists x \phi(x) \rightarrow \exists x \exists \alpha \in \text{Ord} (\phi(x) \wedge \Theta(\alpha, x_{m_\phi}, x)) \text{ and}$$

$$\exists x (\phi(x) \wedge x \in X) \rightarrow \phi(x_{n_\phi}).$$

Here $\Theta(\cdot, \cdot, \cdot)$ is a formula defining a well-order relative to a cble set X for an X -premouse and \dot{X} is the symbol for X in Lpm .

(4) T is a complete, consistent theory s.t. for every model M of T and every model N^* which is the definable closure of $\{x_i : i < \omega^y\}$ in $M \upharpoonright Lpm$, N^* is well-founded and if N denotes the transitive collapse of N^* ,

- (a) N is a \dot{Y} -small X -premouse, where $X = \{x_i : i < \omega^y\} = N \cap \mathbb{R}$,
- (b) $N \models \text{ZF} + \text{"there are no Woodin cardinals"}$,
- (c) $N \models \text{AD}$,
- (d) if $P \not\in N$ and P satisfies (4b), then $P \models \text{AD}$,
- (e) if N is minimal with (b) and (c),

(e) N is III_2^1 -iterable,

If there is a non-determined set of reals in N definable from z and if $Z(z, N)$ is the least such set (in the well-order relative to z defined by $\Theta(\cdot, z, \cdot)$), then $a \oplus b \in Z(z, N)$.

Rmk: Rule (4) can be followed by Player I by playing an appropriate theory as then the model N is uniquely determined since all possible N^* 's are isomorphic.

By our hypothesis this game is determined.

Case 1: Player I has a winning strategy $\bar{\sigma}$ in G .

Let $W \stackrel{\text{definable}}{=} Y \times V_n$ and consider $M_1(\mathbb{R}^W)$.
 ↑
 trans.
 collapse
 ↓
 large enough

Since there is a club of $\mathbb{R}^W \in \mathcal{P}_{\omega_1}(\mathbb{R})$ for such W we can

assume $M_1(\mathbb{R}^W) \cap \mathbb{R} = \mathbb{R}^W$ (by Lemma 1) and

$M_1(\mathbb{R}^W) \setminus S_{\mathbb{R}^W} \not\models AD$ (by hypothesis).

Note that the game G is definable in W and let $\bar{\sigma} \in W$ be st. $\pi(\bar{\sigma}) = \bar{\sigma}$ (i.e. $\bar{\sigma} = \sigma \cap W$). Then $\bar{\sigma}$ is a winning strategy for I in G inside W .

Let $n: \omega \rightarrow \mathbb{R}^W$ be an enumeration of \mathbb{R}^W , $n \in V$.

Consider the following play p of G in V :

- I plays according to his winning strategy $\bar{\sigma}$.
- II plays some real $b \in {}^\omega \omega$ and then n .

Since every $n \in W$ for new, every initial segment of this play is in the domain of $\bar{\sigma}$ ($= \sigma \cap W$).

Hence I's moves are in W as well and z, x_1, x_2, \dots enumerate \mathbb{R}^W .

Let N_p be the \mathbb{T}_2^1 -stable \mathbb{R}^W -premisse obtained from the fact that $\bar{\sigma}$ is a w.s for I.

By comparing N_p with $M_1(\text{TR}^W) \upharpoonright S_{\text{RW}}$ we actually get that N_p is ω_1 -iterable using the minimality property (4d) of N_p .

Let $z = z(z, N_p)$ for $z = \bar{\gamma}(\emptyset) = \bar{\gamma}(\emptyset)$ be the least non-determined set definable from z in N_p .

Consider the Gale-Stewart game $G(z)$ and the following strategy τ for I:

$$a = \tau(b) \quad \text{iff} \quad (z, a) = \bar{\gamma}(b).$$

Note: $\tau \in W$ and τ can be coded by a real, so $\tau \in N_p$.

Then we can show:

Claim: τ is a winning strategy for I in $G(z)$ in N_p .

This contradicts the fact that z is not determined in N_p .

Case 2: Player II has a winning strategy σ in G .

Analogous, asking I to play the theory of the shortest initial segment of $M_1(\text{TR}^W) \upharpoonright S_{\text{RW}}$ satisfying $ZF + \neg AD$ (for W as in Case 1), together with some real $a_0 \in {}^{\omega_1}\omega$ and an enumeration η of TR^W , while II plays according to his winning strategy.

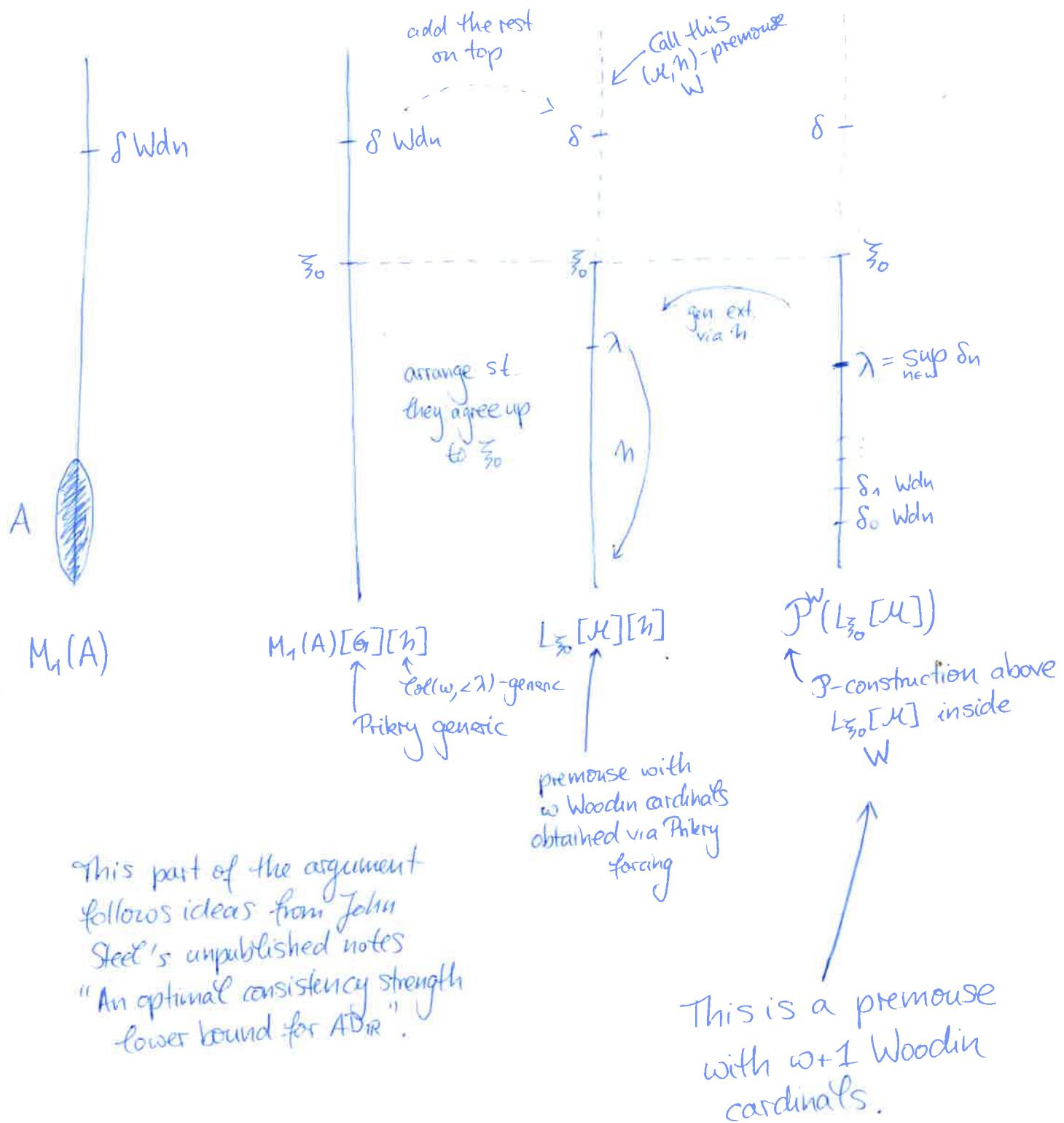
□

Lemma 3: Sp. $A \in \mathcal{P}_{\omega_1}(\text{TR})$ with $M_1(A) \models ZF + AD$ and $M_1(A) \upharpoonright \text{TR} = A$. Then $M_1(A) \models DC$ and moreover we can sp. $M_1(A) \models AD^+ + \Theta = \Theta_0$.

Let's fix an $A \in \mathcal{P}_{\omega_1}(\text{TR})$ as in Lemma 3.

Picture for the rest of the proof.

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A detailed preprint containing a full proof of this result will be on my webpage soon.
Check <https://muelersandra.github.io/publications>.